

Shepherdson's theorems for fragments of open induction

Petr Glivický (joint work with Jana Glivická)

petrglivicky@gmail.com

University of Economics, Prague

Prague Bounded Arithmetic Workshop

November 3, 2017

Slides available at <http://www.glivicky.cz>

Open induction

The **open induction arithmetic** (**IOpen**) is the formal arithmetical theory in the language $L = \langle 0, 1, +, -, \cdot, \leq \rangle$ axiomatized by a set of elementary properties of $0, 1, +, -, \cdot, \leq$ and by the scheme of axioms of open induction:

$$(\varphi(0, \bar{y}) \ \& \ (\forall x)(\varphi(x, \bar{y}) \rightarrow \varphi(x + 1, \bar{y}))) \rightarrow (\forall x)\varphi(x, \bar{y}),$$

for all **open** (i.e. quantifier free) formulas φ of L .

The axioms ensure that any model \mathcal{M}^+ of IOpen is the nonnegative part of the unique **discretely ordered ring** \mathcal{M} .

We often say that a discretely ordered ring \mathcal{M} is a model of IOpen, by which we mean that its nonnegative part \mathcal{M}^+ is.

The Shepherdson's theorem

Question: Is there an algebraic characterisation of those discretely ordered rings that are models of IOpen ?

Theorem (J. C. Shepherdson, 1964)

A discretely ordered ring is a model of IOpen if and only if it is an integer part of its real closure.

The **real closure** $R(\mathcal{M})$ of a discretely ordered ring \mathcal{M} is the unique (ordered) **real closed field** extending the ordered **fraction field** $F(\mathcal{M})$ of \mathcal{M} .

A discretely ordered subring I of \mathcal{R} is called an **integer part** of \mathcal{R} (denoted by $I \subseteq^{IP} \mathcal{R}$) if for every $r \in R$, there is $i \in I$ such that $r - 1 < i \leq r$. Such an i is called the **integer part** of r (w.r.t. I).

Integer-parts-of-roots property

Let $f(x)$ be a definable unary function on \mathcal{M} . By $\text{IPR}(f)$ (*integer-parts-of-roots*) we denote the following formula:

$$(a < b \ \& \ f(a) \leq y < f(b)) \rightarrow (\exists x)(a \leq x < b \ \& \ f(x) \leq y < f(x+1)).$$

The intended meaning of $\mathcal{M} \models \text{IPR}(f)$ can be expressed in vague terms as “existence of integer parts for all f -roots of values $y \in M$ ”.

If \mathcal{F} is a set of definable unary functions on \mathcal{M} , we write $\text{IPR}(\mathcal{F})$ for the scheme $\{\text{IPR}(f); f \in \mathcal{F}\}$.

The Shepherdson's theorem reformulated

We can stratify Shepherdson's theorem a little bit more:

Theorem (J. C. Shepherdson, 1964)

The following are equivalent for any discretely ordered ring \mathcal{M} :

- 1 $\mathcal{M} \subseteq^{IP} R(\mathcal{M})$,
- 2 all roots $r \in R(\mathcal{M})$ of polynomials $p \in \mathcal{M}[x]$ have integer parts in \mathcal{M} ,
- 3 $\mathcal{M} \models \text{IPR}(\mathcal{M}[x])$,
- 4 $\mathcal{M}^+ \models \text{IOpen}$.

Presburger arithmetic

The **Presburger arithmetic** (Pr) is the formal arithmetical theory in the language $L^+ = \langle 0, 1, +, -, \leq \rangle$ axiomatized by a set of elementary properties of $0, 1, +, -, \leq$ and by the scheme of axioms of induction:

$$(\varphi(0, \bar{y}) \ \& \ (\forall x)(\varphi(x, \bar{y}) \rightarrow \varphi(x + 1, \bar{y}))) \rightarrow (\forall x)\varphi(x, \bar{y}),$$

for **all** formulas φ of L^+ .

Every model of Pr is the nonnegative part of the unique **discretely ordered additive Abelian group**.

“Shepherdson’s theorem” for Presburger arithmetic

Theorem

Let \mathcal{M} be a discretely ordered additive Abelian group. Then the following are equivalent:

- 1 $\mathcal{M} \subseteq \text{IP} \frac{M}{\mathbb{N}}$,
- 2 all fractions m/n with $m \in M$ and $0 \neq n \in \mathbb{N}$ have integer parts in M ,
- 3 $\mathcal{M} \models \text{IPR}(\{n(x); n \in \mathbb{N}\})$, where $n(x) = x + \dots + x$ with n summands,
- 4 $\mathcal{M}^+ \models \text{Pr}$.

The Shepherdson's theorem reformulated

Let us recall the Shepherdson's theorem for IOpen:

Theorem (J. C. Shepherdson, 1964)

The following are equivalent for any discretely ordered ring \mathcal{M} :

- 1 $\mathcal{M} \subseteq^{IP} R(\mathcal{M})$,
- 2 all roots $r \in R(\mathcal{M})$ of polynomials $p \in \mathcal{M}[x]$ have integer parts in \mathcal{M} ,
- 3 $\mathcal{M} \models \text{IPR}(\mathcal{M}[x])$,
- 4 $\mathcal{M}^+ \models \text{IOpen}$.

The property $\mathcal{M} \subseteq^{IP} R(\mathcal{M})$ is equivalent to

$$\mathcal{M} \subseteq^{IP} F(\mathcal{M}) \& F(\mathcal{M}) \subseteq^d R(\mathcal{M}),$$

where $R' \subseteq^d R$ denotes that R' is a **dense subring** of R , i.e. for every $q < r$ from R there is $r' \in R'$ such that $q < r' < r$.

Open linear induction

The theory corresponding (in a Shepherdson-like way) to the property $\mathcal{M} \subseteq^{IP} F(\mathcal{M})$ is the the **open linear induction arithmetic**.

The **open linear induction arithmetic** (**IOpenLin**) is the formal arithmetical theory in the language $L = \langle 0, 1, +, -, \cdot, \leq \rangle$ axiomatized by a set of elementary properties of $0, 1, +, -, \cdot, \leq$ and by the scheme of axioms of open linear induction:

$$(\varphi(0, \bar{y}) \ \& \ (\forall x)(\varphi(x, \bar{y}) \rightarrow \varphi(x + 1, \bar{y}))) \rightarrow (\forall x)\varphi(x, \bar{y}),$$

for all **open** (i.e. quantifier free) linear formulas φ of L .

The formula $\varphi(x, \bar{y})$ is called **linear** if in every occurrence of \cdot in φ at least one of the two factors is y_i , for some i .

Theorem

Let \mathcal{M} be a discretely ordered ring. Then the following are equivalent:

- 1 $\mathcal{M} \subseteq^{IP} F(\mathcal{M})$,
- 2 all fractions m'/m with $m', m \in M$ and $0 < m$ have integer parts in M ,
- 3 $\mathcal{M} \models \text{IPR}(\{m(x); m \in M\})$, where $m(x) = m \cdot x$,
- 4 $\mathcal{M} \models (\forall n, k \neq 0)(\exists n')(\exists 0 \leq l < k)(n = n' \cdot k + l)$,
- 5 $\mathcal{M}^+ \models \text{IOpenLin}$.

Open eventual induction

There is a theory that corresponds to the property $F(M) \subseteq^d R(\mathcal{M})$ as well. However, it is not given by a scheme of induction, but by certain generalization of it - the scheme of **open eventual induction**.

The **l -step induction axiom** for the formula $\varphi(x, \bar{y})$ is the following formula denoted by $I'_x\varphi(x, \bar{y})$ or simply $I'_x\varphi$:

$$((\forall u < l)\varphi(u, \bar{y}) \wedge (\forall v)(\varphi(v, \bar{y}) \rightarrow \varphi(v + l, \bar{y}))) \rightarrow (\forall x)\varphi(x, \bar{y}).$$

The **open eventual induction arithmetic** (**eIOpen**) is the formal arithmetical theory in the language $L = \langle 0, 1, +, -, \cdot, \leq \rangle$ axiomatized by a set of elementary properties of $0, 1, +, -, \cdot, \leq$ and by the scheme of axioms of open eventual induction:

$$(\exists l > 0)(I'_x\varphi(x/l, \bar{y})),$$

for all **open** (i.e. quantifier free) formulas φ of L .

Theorem

Let \mathcal{M} be a discretely ordered ring. Then the following are equivalent:

- 1 $F(\mathcal{M}) \subseteq^d R(\mathcal{M})$,
- 2 $(\forall r \in R(\mathcal{M}))(\exists f \in F(\mathcal{M}))(r < f < r + 1)$,
- 3 $(\forall r \in R(\mathcal{M}))(\exists l, m \in M)(m \leq lr < m + l)$,
- 4 $\mathcal{M}^+ \models \text{eIOpen}$.

Alternative axiomatization of IOpen

As we know that

$$(\mathcal{M} \subseteq^{IP} F(\mathcal{M}) \ \& \ F(\mathcal{M}) \subseteq^d R(\mathcal{M})) \rightarrow \mathcal{M} \subseteq^{IP} R(\mathcal{M}),$$

we get the following alternative axiomatization of IOpen:

Corollary

$\text{IOpen} = \text{IOpenLin} + \text{eIOpen}.$

Note that similarly from the inclusion

$$\text{IOpen} \supseteq \text{IOpenLin} + \text{eIOpen},$$

we get:

Corollary

$(\mathcal{M} \subseteq^{IP} F(\mathcal{M}) \ \& \ F(\mathcal{M}) \subseteq^d R(\mathcal{M})) \leftrightarrow \mathcal{M} \subseteq^{IP} R(\mathcal{M}).$

Open questions:

Is there a truly general Shepherdson's theorem, i.e. a theorem such that all variants of Shepherdson theorems for various theories are its special cases?

Are there variants of the Shepherdson's theorem for theories stronger than IOpen ?

Thank you for your attention.

References: [J. Glivická and P. Glivický, *Shepherdson's theorems for fragments of open induction*, APLIMAT 2017 proceedings, 583–589, arXiv: 1701.02001.]